

# On the Computational Complexity of Naive-based Semantics for Abstract Dialectical Frameworks

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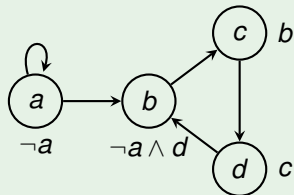
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# Motivation

- **Abstract dialectical frameworks (ADFs):** a powerful generalisation of Dung's abstract argumentation frameworks
- ADFs allow to model argumentation scenarios s.t. that ADF semantics provide interpretations of scenarios
- **Naive-based semantics** are built upon the fundamental concept of conflict-freeness
- **Exhaustive analysis of computational complexity of naive-based semantics**
- Interesting results: some involve little-known classes of so-called **Boolean hierarchy** (another hierarchy in between classes of polynomial hierarchy)
- **Credulous and sceptical entailment:** different complexity depending on whether we check for **truth** or **falsity** of statement

# ADFs and Semantics

## Example



$$\begin{aligned} cfi(D) = \{ & \{a \mapsto u, b \mapsto u, c \mapsto u, d \mapsto u\}, \\ & \{a \mapsto u, b \mapsto t, c \mapsto u, d \mapsto u\}, \\ & \{a \mapsto u, b \mapsto u, c \mapsto t, d \mapsto u\}, \\ & \{a \mapsto u, b \mapsto u, c \mapsto u, d \mapsto t\}, \\ & \{a \mapsto u, b \mapsto t, c \mapsto t, d \mapsto u\}, \\ & \{a \mapsto u, b \mapsto t, c \mapsto u, d \mapsto t\}, \\ & \{a \mapsto u, b \mapsto u, c \mapsto t, d \mapsto t\}, \\ & \{a \mapsto u, b \mapsto t, c \mapsto t, d \mapsto t\}, \\ & \{a \mapsto u, b \mapsto f, c \mapsto f, d \mapsto f\} \end{aligned}$$

$$\begin{aligned} nai(D) = \{ & \{a \mapsto u, b \mapsto t, c \mapsto t, d \mapsto t\}, \\ & \{a \mapsto u, b \mapsto f, c \mapsto f, d \mapsto f\} \end{aligned}$$

## Definition

Let  $D = (S, L, C)$  be an ADF. A three-valued interpretation  $v : S \rightarrow \{t, f, u\}$  is

- **conflict-free**, i.e.  $v \in cfi(D)$ , iff for all  $s \in S$  we have:
  - ▶  $v(s) = t$  implies that  $\varphi_s^v$  is satisfiable,
  - ▶  $v(s) = f$  implies that  $\varphi_s^v$  is unsatisfiable;
- **naive**, i.e.  $v \in nai(D)$ , iff  $v$  is  $\leq_i$ -maximal conflict-free;

# Main Results

	<i>cfi</i>	<i>nai</i>	<i>stg</i>	<i>nai<sub>2</sub></i>
$Ver_\sigma$	DP-c	$\Pi_2^P$ -c	$\Pi_2^P$ -c	$\Pi_2^P$ -c
$Exists_\sigma$	coDP-c	coDP-c	coDP-c	coDP-c
$Cred_\sigma^t$	NP-c	NP-c	$\Sigma_3^P$ -c	$\Sigma_3^P$ -c
$Cred_\sigma^f$	$\Sigma_2^P$ -c	$\Sigma_2^P$ -c	$\Sigma_3^P$ -c	$\Sigma_3^P$ -c
$Scep_\sigma^t$	trivial	$\Pi_2^P$ -c	$\Pi_3^P$ -c	$\Pi_3^P$ -c
$Scep_\sigma^f$	trivial	$\Pi_3^P$ -c	$\Pi_3^P$ -c	$\Pi_3^P$ -c

## Decision Problems

Let  $D = (S, L, C)$  be an ADF,  $\sigma$  a semantics.

- $Ver_\sigma$ : given  $v : S \rightarrow \{t, f, u\}$ , is  $v \in \sigma(D)$ ?
- $Exists_\sigma$ : does there exist a non-trivial interpretation  $v \in \sigma(D)$ ?
- $Cred_\sigma^t / Cred_\sigma^f$ :  $s \in S$ , does there exist an interpretation  $v \in \sigma(D)$  with  $v(s) = t / v(s) = f$ ?
- $Scep_\sigma^t / Scep_\sigma^f$ :  $s \in S$ , is  $v(s) = t / v(s) = f$  for all  $v \in \sigma(D)$ ?