

Complexity Theory  
**Exercise 7: Diagonalization**  
12 December 2017

**Exercise 7.1.** Show the following.

1.  $\text{TIME}(2^n) = \text{TIME}(2^{n+1})$
2.  $\text{TIME}(2^n) \subset \text{TIME}(2^{2n})$
3.  $\text{DTIME}(n) \subset \text{PSPACE}$

**Exercise 7.2.** Find the fault in the following proof of  $P \neq \text{NP}$ :

Assume that  $P = \text{NP}$ . Then  $\text{SAT} \in P$  and thus there exists a  $k \in \mathbb{N}$  such that  $\text{SAT} \in \text{DTIME}(n^k)$ . Because every language in  $\text{NP}$  is polynomial-time reducible to  $\text{SAT}$  we have  $\text{NP} \subseteq \text{DTIME}(n^k)$ . It follows that  $P \subseteq \text{DTIME}(n^k)$ . But by the Time Hierarchy Theorem there exist languages in  $\text{DTIME}(n^{k+1})$  that are not in  $\text{DTIME}(n^k)$ , contradicting  $P \subseteq \text{DTIME}(n^k)$ . Therefore,  $P \neq \text{NP}$ .

**Exercise 7.3.** Consider the function  $\text{pad}: \Sigma^* \times \mathbb{N} \rightarrow \Sigma^* \#^*$  defined as  $\text{pad}(s, \ell) = s\#^j$ , where  $j = \max(0, \ell - |s|)$ . In other words,  $\text{pad}(s, \ell)$  adds enough copies of  $\#$  to the end of  $s$  so that the length is at least  $\ell$ .

For some language  $\mathbf{A} \subseteq \Sigma^*$  and  $f: \mathbb{N} \rightarrow \mathbb{N}$  define  $\text{pad}(\mathbf{A}, f) = \{ \text{pad}(s, f(|s|)) \mid s \in \mathbf{A} \}$ .

1. Show that if  $\mathbf{A} \in \text{DTIME}(n^6)$ , then  $\text{pad}(\mathbf{A}, n^2) \in \text{DTIME}(n^3)$ .
2. Show that if  $\text{NEXPTIME} \neq \text{EXPTIME}$ , then  $P \neq \text{NP}$ .
3. Show for every  $\mathbf{A} \subseteq \Sigma^*$  and every  $k \in \mathbb{N}$  that  $\mathbf{A} \in P$  if and only if  $\text{pad}(\mathbf{A}, n^k) \in P$ .
4. Show that  $P \neq \text{DSpace}(n)$ .
5. Show that  $\text{NP} \neq \text{DSpace}(n)$ .

**Exercise 7.4.** Show that there exists a function that is not time-constructible.