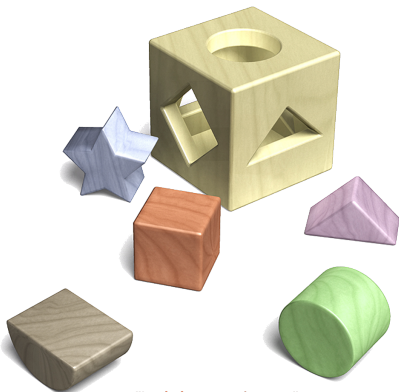


Actions and Causality

Steffen Hölldobler

International Center for Computational Logic
Technische Universität Dresden
Germany

- ▶ Introduction
- ▶ Conjunctive Planning Problems
- ▶ The Fluent Calculus



"Logic is everywhere ..."



States, Actions, and Causality

- ▶ Rational Agents, Agent Programming Languages, Cognitive Robotics
- ▶ **Situation Calculus** McCarthy 1963
- ▶ **Core Ideas**
 - ▷ A state is a snapshot of the world and
 - ▷ can only be changed by actions
- ▶ A state is specified with the help of fluents
- ▶ **Problem** Each state and each action is only partially known!



General Problems

- ▶ **Frame problem**
Which fluents are unaffected by the execution of an action?
- ▶ **Ramification problem**
Which fluents are really present after the execution of an action?
- ▶ **Qualification problem**
Which preconditions have to be satisfied such that an action is executable?
- ▶ **Prediction problem**
How long are fluents present in certain situations?

- ▶ All problems have a cognitive as well as a technical aspect
- ▶ Only the frame problem is considered in this lecture



Requirements

- ▶ **McCarthy 1963**
- ▶ **General properties of causality and facts about the possibility and results of actions are given as formulas**
- ▶ **It is a logical consequence of the facts of a state and the general axioms that goals can be achieved by performing certain actions**
- ▶ **The formal descriptions of states should correspond as closely as possible to what people may reasonably be presumed to know about them when deciding what to do**



Conjunctive Planning Problems

- ▶ **Initial state** $\mathcal{I} : \{i_1, \dots, i_m\}$ of ground fluents
- ▶ **Goal state** $\mathcal{G} : \{g_1, \dots, g_n\}$ of ground fluents
- ▶ Finite set \mathcal{A} of **actions** of the form

$$\{c_1, \dots, c_l\} \Rightarrow \{e_1, \dots, e_k\},$$

where $\{c_1, \dots, c_l\}$ and $\{e_1, \dots, e_k\}$ are multisets of fluents called **conditions** and **(direct) effects**, respectively

- ▶ **Assumption**
Each variable occurring in the effects of an action occurs also in its conditions
- ▶ A **conjunctive planning problem** is the question of whether there exists a sequence of actions whose execution transforms the initial into the goal state



Actions and Plans

- ▶ Let \mathcal{S} be a multiset of ground fluents
- ▶ $\mathcal{C} \Rightarrow \mathcal{E}$ is **applicable** in \mathcal{S} **iff** there exists θ such that $\mathcal{C}\theta \dot{\subseteq} \mathcal{S}$
- ▶ The **application** of $\mathcal{C} \Rightarrow \mathcal{E}$ in \mathcal{S} leads to $\mathcal{S}' = (\mathcal{S} \setminus \mathcal{C}\theta) \dot{\cup} \mathcal{E}\theta$
 - ▷ One should observe that \mathcal{S}' is ground
 - ▶▶ \mathcal{S} is ground
 - ▶▶ $\text{var}(\mathcal{E}) \subseteq \text{var}(\mathcal{C})$
 - ▶▶ θ is grounding
- ▶ A **plan** is a sequence of actions
- ▶ A goal \mathcal{G} is **satisfied** **iff** there exists a plan p which transforms \mathcal{I} into \mathcal{S} and $\mathcal{G} \dot{\subseteq} \mathcal{S}$
 - ▷ Such a plan is called **solution** for the planning problem



Blocks World

▶ **The pickup action**

$$\text{pickup}(V) : \{ \text{clear}(V), \text{ontable}(V), \text{empty} \} \Rightarrow \{ \text{holding}(V) \}$$

▶ **The unstack action**

$$\text{unstack}(V, W) : \{ \text{clear}(V), \text{on}(V, W), \text{empty} \} \Rightarrow \{ \text{holding}(V), \text{clear}(W) \}$$

▶ **The putdown action**

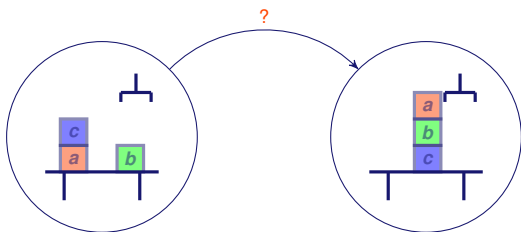
$$\text{putdown}(V) : \{ \text{holding}(V) \} \Rightarrow \{ \text{clear}(V), \text{ontable}(V), \text{empty} \}$$

▶ **The stack action**

$$\text{stack}(V, W) : \{ \text{holding}(V), \text{clear}(W) \} \Rightarrow \{ \text{on}(V, W), \text{clear}(V), \text{empty} \}$$



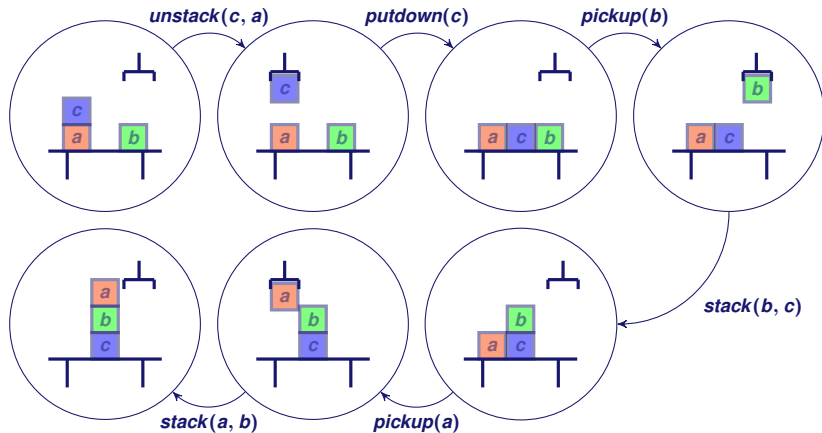
Sussman's Anomaly



- ▶ $\mathcal{I} = \{ \text{ontable}(a), \text{ontable}(b), \text{on}(c, a), \text{clear}(b), \text{clear}(c), \text{empty} \}$
- ▶ $\mathcal{G} = \{ \text{ontable}(c), \text{on}(b, c), \text{on}(a, b), \text{clear}(a), \text{empty} \}$
- ▶ **Solution**
 $[\text{unstack}(c, a), \text{putdown}(c), \text{pickup}(b), \text{stack}(b, c), \text{pickup}(a), \text{stack}(a, b)]$
- ▶ What happens if we independently search for shortest solutions for the two subgoals $\text{on}(a, b)$ and $\text{on}(b, c)$?



Sussman's Anomaly – Solution



A Fluent Calculus Implementation – Actions and Causality

- ▶ An action $C \Rightarrow \mathcal{E}$ is represented by *action*(C^{-l} , *name*, \mathcal{E}^{-l}), where *name* is a term identifying the action

action(*clear*(V) \circ *ontable*(V) \circ *empty*, *pickup*(V), *holding*(V))
action(*clear*(V) \circ *on*(V, W) \circ *empty*, *unstack*(V, W), *holding*(V) \circ *clear*(W))
action(*holding*(V), *putdown*(V), *clear*(V) \circ *ontable*(V) \circ *empty*)
action(*holding*(V) \circ *clear*(W), *stack*(V, W), *on*(V, W) \circ *clear*(V) \circ *empty*)

- ▶ Causality is expressed by *causes*(s, p, s'), where s and s' are fluent terms and p is a list of actions representing a plan

causes($X, [], Y$) $\leftarrow X \approx Y \circ Z$
causes($X, [V|W], Y$) \leftarrow *action*(P, V, Q)
 $\wedge P \circ Z \approx X$
 \wedge *causes*($Z \circ Q, W, Y$)

$X \approx X$



A Fluent Calculus Implementation – The Planning Problem

- ▶ Let \mathcal{K}_A be the set of facts representing actions
- ▶ Let \mathcal{K}_C be the set of clauses representing causality
- ▶ The planning problem (with initial and goal state \mathcal{I} and \mathcal{G} , respectively) is the problem whether

$$\mathcal{K}_A \cup \mathcal{K}_C \cup \mathcal{E}_{AC1} \cup \mathcal{E}_{\approx} \models (\exists P) \text{causes}(\mathcal{I}^{-I}, P, \mathcal{G}^{-I})$$

holds



SLDE-Resolution

▶ Let

- ▶ \mathcal{K} be a set of definite clauses not containing \approx in their heads
- ▶ \mathcal{E} be an equational system and
- ▶ G a goal clause

▶ **Question** Does $\mathcal{K} \cup \mathcal{E} \cup \mathcal{E}_{\approx} \models \neg \forall G$ hold?

- ▶ Let C be a new variant $H \leftarrow A_1 \wedge \dots \wedge A_m$ of a clause in $\mathcal{K} \cup \{X \approx X\}$, G the goal clause $\leftarrow B_1 \wedge \dots \wedge B_n$, and $UP_{\mathcal{E}}$ an \mathcal{E} -unification procedure. If H and B_i , $i \in [1, n]$, are \mathcal{E} -unifiable with $\theta \in UP_{\mathcal{E}}(H, B_i)$ then

$$\leftarrow (B_1 \wedge \dots \wedge B_{i-1} \wedge A_1 \wedge \dots \wedge A_m \wedge B_{i+1} \wedge \dots \wedge B_n)\theta$$

is called **SLDE-resolvent** of C and G

▶ **Theorem 4.10**

- ▶ SLDE-resolution is sound if $UP_{\mathcal{E}}$ is sound
- ▶ SLDE-resolution is complete if $UP_{\mathcal{E}}$ is complete
- ▶ The selection of the literal B_i is don't care non-deterministic



A Solution to Sussman's Anomaly (1)

- (1) \leftarrow *causes*(*ontable*(*a*) \circ *ontable*(*b*) \circ *on*(*c*, *a*) \circ *clear*(*b*) \circ *clear*(*c*) \circ *empty*,
 W ,
ontable(*c*) \circ *on*(*b*, *c*) \circ *on*(*a*, *b*) \circ *clear*(*a*) \circ *empty*).
- (2) \leftarrow *action*(P_1, V_1, Q_1) \wedge
 $P_1 \circ Z_1 \approx$ *ontable*(*a*) \circ *ontable*(*b*) \circ *on*(*c*, *a*) \circ *clear*(*b*) \circ *clear*(*c*) \circ *empty* \wedge
causes($Z_1 \circ Q_1, W_1, \textit{ontable}(\textit{c}) \circ \textit{on}(\textit{b}, \textit{c}) \circ \textit{on}(\textit{a}, \textit{b}) \circ \textit{clear}(\textit{a}) \circ \textit{empty}$).
- (3) \leftarrow *clear*(V_2) \circ *on*(V_2, W_2) \circ *empty* \circ $Z_1 \approx$
ontable(*a*) \circ *ontable*(*b*) \circ *on*(*c*, *a*) \circ *clear*(*b*) \circ *clear*(*c*) \circ *empty* \wedge
causes($Z_1 \circ \textit{holding}(V_2) \circ \textit{clear}(W_2),$
 $W_1,$
ontable(*c*) \circ *on*(*b*, *c*) \circ *on*(*a*, *b*) \circ *clear*(*a*) \circ *empty*).
- (4) \leftarrow *causes*(*ontable*(*a*) \circ *ontable*(*b*) \circ *clear*(*b*) \circ *clear*(*a*) \circ *holding*(*c*),
 $W_1,$
ontable(*c*) \circ *on*(*b*, *c*) \circ *on*(*a*, *b*) \circ *clear*(*a*) \circ *empty*).
- ⋮



A Solution to Sussman's Anomaly (2)

- (7) \leftarrow *causes*(*ontable(a)* \circ *ontable(b)* \circ *clear(b)* \circ
clear(a) \circ *clear(c)* \circ *ontable(c)* \circ *empty*, W_4 ,
ontable(c) \circ *on(b, c)* \circ *on(a, b)* \circ *clear(a)* \circ *empty*).
- ⋮
- (10) \leftarrow *causes*(*ontable(a)* \circ *clear(c)* \circ *ontable(c)* \circ *clear(a)* \circ *holding(b)*, W_7 ,
ontable(c) \circ *on(b, c)* \circ *on(a, b)* \circ *clear(a)* \circ *empty*).
- ⋮
- (14) \leftarrow *causes*(*ontable(a)* \circ *ontable(c)* \circ *clear(a)* \circ *on(b, c)* \circ *clear(b)* \circ *empty*, W_{10} ,
ontable(c) \circ *on(b, c)* \circ *on(a, b)* \circ *clear(a)* \circ *empty*).
- ⋮
- (16) \leftarrow *causes*(*ontable(c)* \circ *on(b, c)* \circ *clear(b)* \circ *holding(a)*, W_{13} ,
ontable(c) \circ *on(b, c)* \circ *on(a, b)* \circ *clear(a)* \circ *empty*).
- ⋮
- (19) \leftarrow *causes*(*ontable(c)* \circ *on(b, c)* \circ *clear(a)* \circ *on(a, b)* \circ *empty*, W_{16} ,
ontable(c) \circ *on(b, c)* \circ *on(a, b)* \circ *clear(a)* \circ *empty*).
- (20) []



Solving the Frame Problem

- ▶ In the fluent calculus the frame problem is mapped onto fluent matching and fluent unification problems
- ▶ For example, let

$$s = \text{ontable}(a) \circ \text{ontable}(b) \circ \text{on}(c, a) \circ \text{clear}(b) \circ \text{clear}(c) \circ \text{empty}$$

$$t = \text{clear}(c) \circ \text{on}(c, a) \circ \text{empty}$$

then

$$\theta = \{Z \mapsto \text{ontable}(a) \circ \text{ontable}(b) \circ \text{clear}(b)\}$$

is a most general \mathcal{E} -matcher for the \mathcal{E} -matching problem

$$\mathcal{E}_{AC1} \cup \mathcal{E}_{\approx} \models (\exists Z) s \approx t \circ Z$$

- ▶ Consequently, $\text{unstack}(c, a)$ can be applied to s yielding

$$s' = \text{ontable}(a) \circ \text{ontable}(b) \circ \text{clear}(b) \circ \text{clear}(a) \circ \text{holding}(c)$$



Why are States not Modelled by Sets?

- ▶ Let $\mathcal{E}_{AC1} = \mathcal{E}_{AC1} \cup \{X \circ X \approx X\}$
- ▶ In this case the \mathcal{E} -matching problem

$$\mathcal{E}_{AC1} \cup \mathcal{E}_{\approx} \models (\exists Z) s \approx t \circ Z$$

has an additional solution, viz.

$$\eta = \{Z \mapsto \text{ontable}(a) \circ \text{ontable}(b) \circ \text{clear}(b) \circ \text{empty}\}$$

θ and η are independent wrt \mathcal{E}_{AC1}

- ▶ Computing the successor state in this case yields

$$s'' = \text{ontable}(a) \circ \text{ontable}(b) \circ \text{clear}(b) \circ \text{clear}(a) \circ \text{holding}(c) \circ \text{empty}$$

which is not intended because the arm of the robot cannot be empty and holding an object at the same time



Remarks (1)

- ▶ **Some people even believed that the frame problem cannot be solved within first order logic**
- ▶ **Forward versus backward planning**
- ▶ **Many extensions**
 - ▷ **Incomplete specifications of initial situation, e.g.**

$$\begin{aligned}
 &(\exists X, P, Y) \\
 &\text{causes}(\text{ontable}(b) \circ Y, \\
 &\quad P, \\
 &\quad \text{ontable}(c) \circ \text{on}(b, c) \circ \text{on}(a, b) \circ \text{clear}(a) \circ \text{empty} \circ X)
 \end{aligned}$$

- ▷ **Indeterminate effects**
- ▷ **Specificity**
- ▷ **Ramification and qualification problem**



Remarks (2)

- ▶ **Fluent calculus versus linear logic versus linear connection method**
- ▶ **Fluent calculus versus situation calculus versus event calculus**
- ▶ **Planning problems can be reduced to SAT-problems if the length of a plan is restricted**

