EXERCISE 1 Science of Computational Logic

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International Masters Programme in Computational Logic — winter semester 2015/2016

Problem 1.1

In the lectures the following example from Description Logics was presented:

- $$\begin{split} \mathcal{K}_T : & \text{woman} \sqsubseteq \text{person}, & \mathcal{K}_A : \\ & \text{man} \sqsubseteq \text{person}, \\ & \text{mother} = \text{woman} \sqcap \exists \text{child} : \text{person}, \\ & \text{father} = \text{man} \sqcap \exists \text{child} : \text{person}, \\ & \text{parent} = \text{mother} \sqcup \texttt{father}, \\ & \text{grandparent} = \text{parent} \sqcap \exists \text{child} : \text{parent}, \\ & \text{father_without_son} = \text{father} \sqcap \forall \text{child} : \neg \text{man} \end{split}$$
 - parent(carl), parent(conny), child(conny, joe), child(conny, carl), man(joe), man(carl), woman(conny).

Are the following consequences valid? Justify your answers.

- 1. $\mathcal{K}_T \cup \mathcal{K}_A \models \mathsf{grandparent}(\mathsf{conny})$
- 2. $\mathcal{K}_T \cup \mathcal{K}_A \models \mathsf{father}(\mathsf{carl})$
- 3. $\mathcal{K}_T \cup \mathcal{K}_A \models \mathsf{father_without_son(carl)}$

Problem 1.2

Prove that $F \sqsubseteq G \equiv F \sqcap \neg G = \bot$

Problem 1.3

Show that grandparent $\sqsubseteq_{\mathcal{K}_T}$ parent by reducing subsumption into concept satisfiability, where \mathcal{K}_T is the T-Box from Problem 1.1.

Problem 1.4

Is the concept (father \sqcap mother) satisfiable w.r.t. \mathcal{K}_T from Problem 1.1?

Problem 1.5

1. Which generalized concept axioms must be added to prevent that a person is female and male?

2. Is there a single generalized concept axiom that prevents that a person is female and male?

Problem 1.6

Give an equivalent concept of (woman $\sqcap \exists \mathsf{child.person})$ without using the constructors \sqcap and $\exists r.C$

Problem 1.7

Prove the following:

If $(\forall r.C)(a) \in \mathcal{A}$, and $r(a,b) \in \mathcal{A}$, then $\mathcal{A} \models C(b)$.

Problem 1.8

Prove the following:

If $(\exists r.C)(a) \in \mathcal{A}$, \mathcal{A} is satisfiable, and b is a Skolem constant, then $\mathcal{A} \cup \{r(a,b), C(b)\}$ is satisfiable as well.

Problem 1.9

Let ${\mathcal A}$ be an ABox. Proof or refute the following claims:

- 1. If \mathcal{A} contains only elements of the form r(a,b) where r is a role name and a,b are individual names, then \mathcal{A} is satisfiable.
- 2. If \mathcal{A} contains only elements of the form A(x) where A is a concept name and a is an individual name, then \mathcal{A} is satisfiable.
- 3. If \mathcal{A} contains only elements of the form A(x) or $\neg A(x)$, where A is an atomic concept name and x an individual name, then \mathcal{A} is satisfiable.