Exercise Sheet 12: Dependencies David Carral, Markus Krötzsch Database Theory, July 17, Summer Term 2018

Exercise 12.1. Let \mathcal{L} be a fragment of first-order logic for which finite model entailment and arbitrary model entailment coincide, i.e., for every \mathcal{L} -theory \mathcal{T} and every \mathcal{L} -formula φ , we find that φ is true in all models of \mathcal{T} if and only if φ is true in all finite models of \mathcal{T} .

- (a) Give an example for a proper fragment of first-order logic with this property.
- (b) Give an example for a proper fragment of first-order logic without this property.
- (c) Show that entailment is decidable in any fragment with this property.

Exercise 12.2. Consider the following set of tgds Σ :

$$A(x) \to \exists y. R(x, y) \land B(y)$$
$$B(x) \to \exists y. S(x, y) \land A(y)$$
$$R(x, y) \to S(y, x)$$
$$S(x, y) \to R(y, x)$$

Does the oblivious chase universally terminate for Σ ? What about the restricted chase?

Exercise 12.3. Is the following set of tgds weakly acyclic?

$$B(x) \to \exists y. S(x, y) \land A(x)$$
$$A(x) \land C(x) \to \exists y. R(x, y) \land B(y)$$

Does the oblivious chase universally terminate over this set of tgds?

Exercise 12.4. Termination of the oblivious (resp. restricted) chase over a set of tgds Σ implies the existence of a finite universal model for Σ . Is the converse true? That is, does the existence of a finite universal model for Σ imply termination of the oblivious (resp. restricted) chase?

Exercise 12.5. A term is *cyclic* if it is of the form $f(t_1, \ldots, t_n)$ and, for some $i \in \{1, \ldots, n\}$, the function symbol f syntactically occurs in t_i . A set Σ of tgds that does not contain any constants is *model-faithful acyclic* (MFA) iff no cyclic term occurs in the skolem chase of $\Sigma \cup \mathcal{I}_{\star}$ (with \mathcal{I}_{\star} the critical instance). Show the following:

- Checking MFA membership is decidable.
- If Σ is MFA, then the skolem chose universally terminates for Σ .