# PROBLEM SOLVING AND SEARCH IN ARTIFICIAL INTELLIGENCE 

## Lecture 5 Tabu Search

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## Agenda

(1) Introduction
(2) Constraint Satisfaction (CSP)
(3) Uninformed Search versus Informed Search (Best First Search, A* Search, Heuristics)
4) Local Search, Stochastic Hill Climbing, Simulated Annealing
(5) Tabu Search
(6) Answer-set Programming (ASP)
(7) Structural Decomposition Techniques (Tree/Hypertree Decompositions)
(8) Evolutionary Algorithms/ Genetic Algorithms

## Tabu Search

## Main Idea

- A memory forces the search to explore new areas of the search space
- Memorize solutions that have been examined recently. They become tabu points in next steps
- Tabu search is deterministic


## Tabu Search and SAT

- SAT problem with $n=8$ variables
- Initial (random) assignment $\mathbf{x}=(0,1,1,1,0,0,0,1)$
- Evaluation function: weighted sum of number of satisfied clauses. Weights depend on the number of variables in the clause
- Maximize evaluation function (i.e. we're trying to satisfy all clauses)
- Random assignment provides $\operatorname{eval}(\mathbf{x})=27$
- Neighborhood of $\mathbf{x}$ consists of 8 solutions. Evaluate them and select best
- At this stage, it is the same as hill-climbing
- Suppose flipping 3rd variable generates best evaluation $\left(\operatorname{eval}\left(\mathbf{x}^{\prime}\right)=31\right)$
- Memory keeps track of actions


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## Question

(1) What is stored in memory (think of SAT as an example)?
(2) How can we escape local optima with help of the memory?

## Recency-based Memory

- Index of flipped variable + time when it was flipped
- Differentiate between older and more recent flips
- SAT: time stamp for each position of solution vector $M$ (initialized to 0 )
- Value of time stamp provides information on recency of flip at position


## Memory Vector

$$
M(i)=j(\text { when } j \neq 0)
$$

$j$ is most recent iteration when $i$-th bit was flipped

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Assume information is stored for at most 5 iterations.

## Alternative Interpretation

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$i$-th bit was flipped $5-j$ iterations ago

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$$
i \text {-th bit was flipped } 5-j \text { iterations ago }
$$

## Example

| 0 | 0 | 5 | 0 | 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Memory after one iteration. 3rd bit is tabu for next 5 iterations.

## Different Interpretations

## 1st Variant

- Stores iteration number of most recent flip
- Requires a current iteration counter $t$ which is compared with memory values
- If $t-M(i)>5$ forget
- Only requires updating a single entry, and increase the counter
- Used in most implementations


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## 2nd Variant

- Values are interpreted as number of iterations for which a position is not available
- All nonzero entries are decreased by one at every iteration


## Example ctd.

- Initial assignment $\mathbf{x}=(0,1,1,1,0,0,0,1)$
- After 4 additional iterations $M$ :

| 3 | 0 | 1 | 5 | 0 | 4 | 2 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

- Most recent flip $M(4)=5$
- Current solution: $\mathbf{x}=(1,1,0,0,0,1,1,1)$ with $\operatorname{eval}(\mathbf{x})=33$


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## Neighborhood of $\mathbf{x}$

$$
\begin{array}{ll}
\mathbf{x}_{1}=(0,1,0,0,0,1,1,1) & \mathbf{x}_{5}=(1,1,0,0,1,1,1,1) \\
\mathbf{x}_{2}=(1,0,0,0,0,1,1,1) & \mathbf{x}_{6}=(1,1,0,0,0,0,1,1) \\
\mathbf{x}_{3}=(1,1,1,0,0,1,1,1) & \mathbf{x}_{7}=(1,1,0,0,0,1,0,1) \\
\mathbf{x}_{4}=(1,1,0,1,0,1,1,1) & \mathbf{x}_{8}=(1,1,0,0,0,1,1,0)
\end{array}
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## Neighborhood of $x$

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\mathbf{x}_{4}=(1,1,0,1,0,1,1,1) & \mathbf{x}_{8}=(1,1,0,0,0,1,1,0)
\end{array}
$$

TABU, best evaluation $\operatorname{eval}\left(\mathbf{x}_{5}\right)=32$, decrease!

## Example ctd.

- Current solution: $\mathbf{x}=(1,1,0,0,0,1,1,1)$ with $\operatorname{eval}(\mathbf{x})=33$
- New solution: $\mathbf{x}_{5}=(1,1,0,0,1,1,1,1)$ with $\operatorname{eval}\left(\mathbf{x}_{5}\right)=32$

| 3 | 0 | 1 | 5 | 0 | 4 | 2 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

changes to:

| 2 | 0 | 0 | 4 | 5 | 3 | 1 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

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changes to:

| 2 | 0 | 0 | 4 | 5 | 3 | 1 | 0 |
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## Policy might be too restrictive

- What if tabu neighbor $\mathbf{x}_{6}$ provides excellent evaluation score?
- Make search more flexible: override tabu classification if solution is outstanding
$\Longrightarrow$ aspiration criterion


## Frequency-based Memory

- Operates over a longer horizon
- SAT: vector $H$ serves as long-term memory.
- Initialized to 0, at any stage of the search

$$
H(i)=j
$$

interpreted as: during last $h$ (horizon) iterations, the $i$-th bit was flipped $j$ times

- Usually horizon is large
- After 100 iterations with $h=50$, long-term memory $H$ might have the following values

| 5 | 7 | 11 | 3 | 9 | 8 | 1 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

- Shows distribution of moves throughout the last 50 iterations


## Diversity of Search

Frequency-based memory provides information about which flips have been under-represented or not represented.
$\Longrightarrow$ we can diversify the search by exploring these possibilities

## Use of Long-term Memory

## Special Circumstances

- Situations where all non-tabu moves lead to worse solution
- To make a meaningful decision about which direction to explore next
- Typically: most frequent moves are less attractive
- Value of evaluation score is decreased by some penalty measure that depends on frequency, final score implies the winner


## Example SAT

- Assume value of current solution is $\operatorname{eval}(\mathbf{x})=35$
- Non-tabu flips 2, 3 and 7 have values 30,33,31
- None of tabu moves provides value greater than 37 (highest value so far) $\Longrightarrow$ we can't apply aspiration criterion


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- Frequency based-memory and evaluation function for new solution $\mathbf{x}^{\prime}$ is

$$
\operatorname{eval}\left(\mathbf{x}^{\prime}\right)-\operatorname{penalty}\left(\mathbf{x}^{\prime}\right)
$$

- penalty $\left(\mathbf{x}^{\prime}\right)=0.7 \times H(i)$, where 0.7 coefficient, $H(i)$ value from long-term memory $H$ :

7
11
1
for solution created by flipping 2nd bit for solution created by flipping 3nd bit for solution created by flipping 7nd bit

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7 for solution created by flipping 2nd bit
11 for solution created by flipping 3nd bit
1 for solution created by flipping 7nd bit

- New scores are:

$$
\begin{array}{lc}
30-0.7 \times 7=25.1 & \text { 2nd bit } \\
33-0.7 \times 11=25.3 & \text { 3nd bit } \\
31-0.7 \times 1=30.3 & \text { 7th bit }
\end{array}
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11 for solution created by flipping 3nd bit
1 for solution created by flipping 7nd bit

- New scores are:

| $30-0.7 \times 7=25.1$ | 2nd bit |
| :--- | ---: |
| $33-0.7 \times 11=25.3$ | 3nd bit |
| $31-0.7 \times 1=30.3$ | 7th bit |

## Diversify Search

Including frequency values in a penalty measure for evaluating solutions.

## Further Options to Diversify Search

We migth add additional rules:

- Aspiration by default: select the oldest of all considered
- Aspiration by search direction: memorize whether or not the performed moves generated any improvement
- Aspiration by influence: measures the degree of change of the new solution
a) in terms of the distance between old and new solution
b) change in solution's feasibility, if we deal with a constraint problem
- Intuition: particular move has a larger influence if a larger step was made from old to new solution


## Summary

- Simulated annealing and tabu search are both design to escape local optima
- Tabu search makes uphill moves only when it is stuck in local optima
- Simulated annealing can make uphill moves at any time
- Simulated annealing is stochastic, tabu search is deterministic
- Compared to classic algorithms, both work on complete solutions. One can halt them at any iteration and obtain a possible solution
- Both have many parameters to worry about


## References

Z. Zbigniew Michalewicz and David B. Fogel. How to Solve It: Modern Heuristics, volume 2. Springer, 2004.

